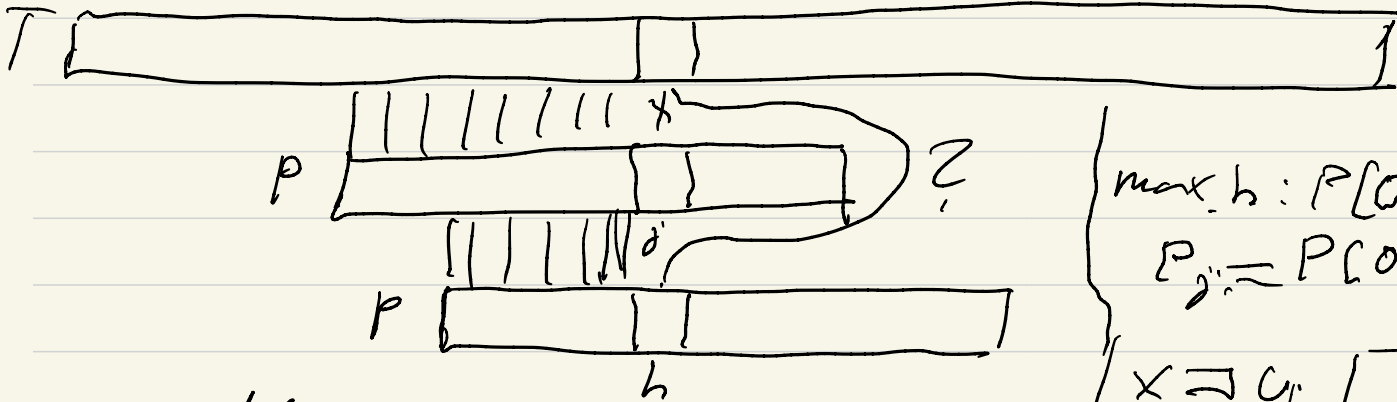


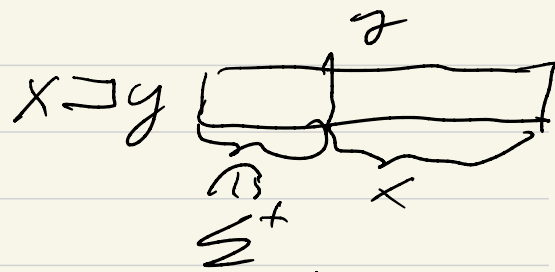
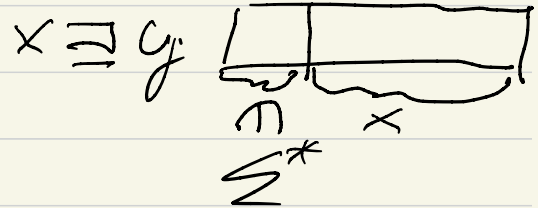
KMP algorithm

$T: \Sigma^n$; $P: \Sigma^m$



$$\max_h: P[0..h] \supseteq P[0..j]$$

$$P_j := P[0..j] \parallel P_h \supseteq P_j$$



$$x \supseteq y \parallel x \sqsubseteq y$$

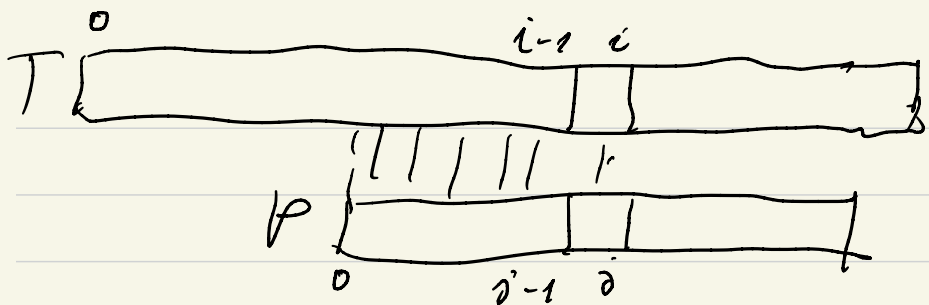
$$D_1 H(j) \stackrel{\text{def}}{=} \{ b \in 0..j-1 \mid P_b \supseteq P_j \}$$

$$x \sqsupseteq y \stackrel{\text{def}}{\iff} x \sqsubseteq y \wedge x \supseteq y \quad (j' \in \{1..m\})$$

$$H(j) = \{ |x| \mid x \sqsupseteq P_j \}$$

D, next: $1..m \rightarrow 0..m-1$

$$\text{next}(j) = \max H(j)$$



$$P_{j+1} \equiv T_{i+1}$$

$$P_j \equiv T_i \wedge P[j] = T[i]$$

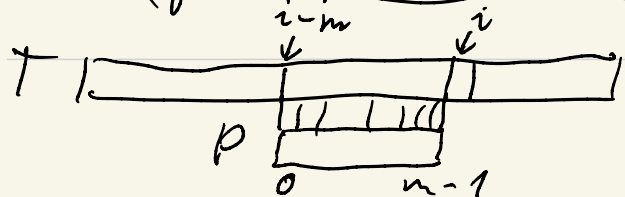
Tac

$$P_{j+1} \equiv T_{i+1} \Leftrightarrow P_j \wedge T_i \wedge P[j] = T[i]$$

$$\text{next}(j) \in 0..j-1 \quad (j \in 1..m)$$

$$\text{next}(j+1) \leq \text{next}(j) + 1 \quad (j \in 1..m-1)$$

$P[j-1] =$	B	A	B	A	B	B	A	B
$j =$	1	2	3	4	5	6	7	8
$\text{next}(j) =$	0	0	1	2	3	1	2	3



$KMP(T: \Sigma^n; P: \Sigma^m; S: \mathcal{N}^{\{ \}})$

```

next[1: N[m]; init(next, P)
S := { }; i := j := 0
while i < n
  P[j] = T[i]
  i++; j++
  if j = m
    S := S ∪ {i-m}
    j := next[m]
  else if j > 0
    j := next[j]
  else
    j := 0

```

$\exists \text{ next } P_j \Rightarrow T_i$
 $1 \leq j \leq i \leq m$
 $0 \leq j < m$

$t(i, j) := 2i - j \in 0..2m$

Initially: $t(i, 0) = 0$
 with each iteration: $t(i, j)$ strictly increases

} \rightarrow KMP: max. $2n$ iteration
 } KMP: min. n iteration
 (initially: $i=0$ finally: $i=n$
 i increases max. by 1 in each iteration)

$$nT(n, m), MT(n, m) \in \Theta(n) + \Theta(m)$$

$$nT_{init}(n), MT_{init}(m) \in \Theta(m)$$

$$m \leq n$$

$$nT(n, m), MT(n, m) \in \Theta(n)$$

KMP($T : \Sigma[n] ; P : \Sigma[m] ; S : \mathbb{N}\{\}$)

$next/1 : \mathbb{N}[m] ; \text{init}(next, P)$			
$S := \{\} ; i := j := 0$			
$i < n$			
$P[j] = T[i]$			
$i++ ; j++$		$j = 0$	
$j = m$		$i++$	$j := next[j]$
$S := S \cup \{i - m\}$	SKIP		
$j := next[m]$			

$$\text{next}(1) = 0 \quad (1 \leq j < m)$$

$$\begin{aligned} \text{next}(j+1) &= \max H(j+1) = \\ &= \max \{h \in 0..j \mid P_h \supseteq P_{j+1}\} = \quad // i := h-1 \\ &= \max \{i+1 \in 0..j \mid P_{i+1} \supseteq P_{j+1}\} = \\ &= \max(\{i+1 \in 1..j \mid P_{i+1} \supseteq P_{j+1}\} \cup \{0\}) = \\ &= \max(\{i+1 \in 1..j \mid \underbrace{P_i \supseteq P_j}_{i \in H(j)} \wedge P[i] = P[j]\} \cup \{0\}). \end{aligned}$$

$$h(j) := \{i \in 0..j-1 \mid \underbrace{i \in H(j)}_{i \in H(j)} \wedge P[i] = P[j]\}$$

$$h(j) \subseteq H(j)$$

$$\text{next}(j+1) = \begin{cases} \max h(j) + 1 & \text{if } h(j) \neq \{\} \\ 0 & \text{if } h(j) = \{\} \end{cases}$$

$$\overline{\Gamma}_+ \max_{l+1} H(j) = \text{next}(\max_l H(j))$$

$$(\max_l H(j) = \text{next}(j))$$

init(next[1..N][m]; P: E[m])

$j := 1; \text{next}[1] := i := 0$		
$j < m$		
$P[0] = P[j]$		
$i++; j++$	$i = 0$	
$\text{next}[j] := i$	$j++$	$i := \text{next}[i]$
	$\text{next}[j] := 0$	

Invariant

$\text{next}[1..j] = \text{next}(1..j)$

$\wedge 0 \leq i < j \leq m$

$\wedge P_i \supseteq P_j \quad (i \in H(j))$

to prove: $MT_{\text{init}}(m), MT_{\text{init}}(m) \in \Theta(m)$

See in the printed Δ Lecture notes !!!

