Types in CLEAN CLEAN

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2009. május 21.



Outline

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Type aliases

```
:: Name :== String
```

- Documentary nature
 - → improves readablity
- Like typedef in C++

Excercise – Synonyms

Write a new icl file, containing

• type synonyms Length for integer! Create a Start function to test it.



Solution – Synonyms

```
:: Length :== Int
```

Tuples

```
:: Point2D :== (Int,Int)
:: Point3D :== (Int,Int,Int)
```

Using tuples

```
mirror :: Point2D \rightarrow Point2D
mirror (x,y) = (y, -x)
```

Tuples elements of different types

Excercise – Tuples

Write a new icl file, containing

- The previously declared Bag type.
- Write a function with the following signature:

```
multiplicity :: Bag String -> Int returns the multiplicity of the given string or zero.
```

A Start function to test it:

```
Start = multiplicity [("a", 1), ("b",2)] "q"
```

Solution – Tuples

Example

Declaration

Value

```
p1 :: Point
p1 = { x = 1
    , y = 1
}
```

Referencing

```
x1 :: Int
x1 = p1.x
```

Comparison with tuple

```
With tuple
:: Point :== (Int, Int)

norm1 :: Point -> Int
norm1 (x,y) = x + y
```

Records can be 'updated'

```
projectX :: Point -> Point
projectX point = { point & x = 0 }
```

- No need to refer to other fields' values
- Makes the usage of records flexible (fields can be added later without having to rewrite all functions)

Exercise – Records

Define new records for pieces of different shapes:

- Square
- and Circle

Use type Length!



Solution

```
:: Square = { side :: Length }
:: Circle = { radius :: Length }
```

Example – Direction

```
:: Direction = Left | Right

mirror :: Direction -> Direction

mirror Left = Right

mirror Right = Left
```

Example - Point

```
:: Point = Point Int Int
Point :: Point -> Int
norm1 Point x y = x + y
```

```
With tuple
:: Point :== (Int, Int)

norm1 :: Point -> Int
norm1 (x,y) = x + y
```

Example – Nat

```
:: Nat = Succ Nat | Zero
One = Succ Zero
Two = Succ One
pred :: Nat -> Nat
pred (Succ n) = n
double :: Nat -> Nat
// 2* (1+n) = 1+1+(2*n)
double (Succ n) = Succ (Succ (double n))
double Zero = 0
```

Composite matching

Example – BiNat

Parametric

```
:: Tree a = Node a (Tree a) (Tree a) | Leaf
height :: Tree a -> Nat
height Leaf = 1
height Node l r = max (height l) (height r)
```

Notation:

- Type constructor: Tree
- Data constructor: Node

Preorder traversal

Algebraic Types – Excercise

- Write an algebraic type Color which contains 4 different colours!
- Write a function that can calculate the number of the leaves in a given tree!
- Write a postorder traversal on the type Tree and save it a list.

Algebraic Types – Solution

```
:: Color = Black | White | Red | Yellow
length :: Tree a -> Int
length Leaf = 1
length Node l r = (length l) + (length r)
```

Algebraic Types – Solution (cont)

Benefits of unique types

- Normally destructive updates can't be done in Clean, because it violates referential transparency.
- Sometimes a destructive update is needed:
 - user interactions
 - writing files or on screen
 - etc.
- We can't refer to a uniquess type twice, so we can safely update it.
- When a unique variable is updated, it is destroyed and a new one created.
- Memory can be reused.



Input/Output

When we read a file it has a side-effect of changing the current position in the file. The next read from the file will return different value, so it violates referential transparency. If the read function returns a new file (with changed current position) we can make the original file unique.

```
write12
# f = open("test.txt")
# f2 = fwritec 'a' f
= 1
```

Arrays

- Contiguous blocks in memory
- Increases efficiency (constant time element selection)
- Increases risk of run-time errors



Lazy vs. Strict Arrays

Lazy: elements are evaluated only when directly referred at

```
:: LazyArray a :== {a}
```

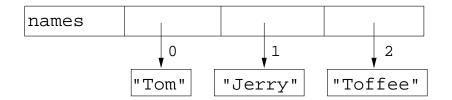
 Strict: elements are evaluated when the array is being referred at

```
:: StrictArray a :== {!a}
```

Boxed arrays

Only pointers to elements are stored in array itself

```
:: BoxedArray a :== {a}
names :: BoxedArray
names = {"Tom", "Jerry", "Toffee"}
```



Unboxed Arrays

Elements are stored directed in array itself

```
:: UnboxedArray a :== {#a}
numbers :: UnboxedArray
numbers = {1, 2, 3, 5, 7, 11}
```

numbers	1	2	3	5	7	11

Array elements

Fields can be referred at directly by their indices (first index is 0):

```
thirdElement array = array.[2]
```

```
Array comprehension
```

```
{ element \\ element <-: array }
```

Unique arrays can be updated:

```
uniqueArray :: *{Int} -> *{Int}
uniqueArray x = { x & [4] = 3, [3] = 4 }
```

Excercise – Arrays

- Create an array EmptyArray, the size of 6 and contains only zeros!
- Write a function, addAmount, which takes an array, the amount of paint needed, and index, that indicates which colour's amount should be modified and updates the array respectively!

Solution – Arrays

```
EmptyArray = { 0.0 \\ i <- [0..NumOfColors-1] }

addAmount array amount index
  # (value, array) = uselect array index
  = { array & [index] = value + amount }

// A bad solution. Two references to array
addAmount array amount index
  = { array & [index] = amount + array.[index] }</pre>
```

Declaration

The type's interface is placed in a separate definition module; user of the type sees only this module

```
stack.dcl
  definition module stack

:: Stack a

Push :: a (Stack a) -> Stack a
  Pop :: (Stack a) -> Stack a
  top :: (Stack a) -> a
  Empty :: Stack a
```

Implementation

Representation of type and implementation of functions is hidden from user of the module (hence the name abstract)

Usage

Representation and implementation can be changed without affecting the modules using this type

```
module stack_user

import stack

Start = top (Push 1 Empty)
```

Abstract Types – Excercise

- Define and implement an abstract type Queue for queueing any types (a queue is a first in first out data-structure)! It has to have the following functions:
 - EmptyQueue
 - insertItem
 - getFirstItem
 - removeFirstItem
- Write a Start expression to test the new queue!

Abstract Types – Solution

```
definition module queue

::Queue a

EmptyQueue :: Queue a
insertItem :: (Queue a) a -> (Queue a)
getFirstItem :: (Queue a) -> a
removeFirstItem :: (Oueue a) -> (Oueue a)
```

Abstract Types – Solution

```
implementation module queue
import StdList
:: Queue a :== [a]

EmptyQueue = []
insertItem queue item = queue ++ [item]
getFirstItem queue = hd queue
removeFirstItem queue = tl queue
```

Abstract Types – Solution (cont)

'Ad hoc' polymorphism: we need the same set of functions for different types, but implementation depends on type

```
pretty printing
class PrettyPrint a
where
   format :: a -> String
   concat :: a a -> String
   concat al a2 = (format al) ++ ",
" ++ (format a2)
```

- We have to instantiate the class for all types we'd like to use it for
- Instantiation can differ from type to type

```
Instatiation for Int
instance PrettyPrint Int
where
    format :: Int -> String
```

format i = "Integer: " ++ toChar i

Instantiation for Point

- Member concat was derived from format
- Usage: user has to indicate that an instantiation of a particular class is needed

Synonyms Tuples Records Algebraic Types Uniqueness typing Arrays Abstract data types Classes

Classes – Excercise

- Write a type class, Measures a that has to functions:
 - circumference
 - surface!
- Instantiate the class to both types of pieces!
- Add a field price to each type of pieces! Does it affect the functions you've just implemented?

Classes – Solution

```
class Measures a
where
    circumference :: a -> Real
    surface :: a -> Real

instance Measures Square
where
    circumference square = fromInt (square.side * 4)
    surface square = fromInt (square.side * square.side)
```

Classes – Solution (cont)